

Finite state verifiers with both private and public coins

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1 Motivation

2 Definitions and notation

3 Results

Motivation

- Previously, by [Con89; GKR15]:

$$\text{IP}(\text{log-space}, \text{poly-public-coins}, \text{poly-time}) = \text{P}$$

- Based on this, our main result:

$$\text{IP}(\text{con-space}, \text{con-private-coins}, \text{poly}^*\text{-public-coins}, \text{poly}^*\text{-time}) = \text{P}$$

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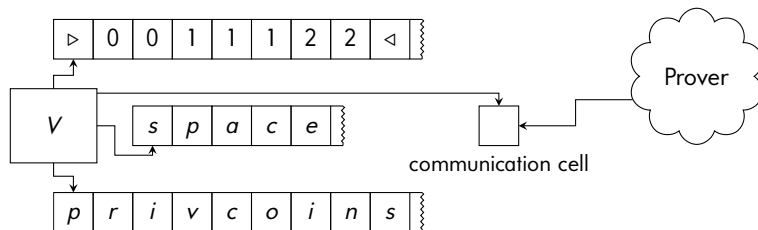
Definitions and notation

■ Interactive Proof System

- Verifier: responsibility and authority to decide membership
- Prover: interactive advocate of membership

■ $IP(\dots\text{-space}, \dots\text{-private-coins}, \dots\text{-public-coins}, \dots\text{-time})$

- Resource budgets of the verifier
- Absent: unavailable



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■ Related examples and **shorthand notations**

- $\text{IP}(\log\text{-space}, \text{poly-public-coins}, \text{poly-time})$

$\mathcal{IP}_{\text{pub-ran}}(\log\text{-space})$

- $\text{IP}(\text{con-space}, \text{con-private-coins}, \text{poly}^*\text{-public-coins}, \text{poly}^*\text{-time})$

$\mathcal{IP}_{\text{pub-ran}}(\text{con-space}, \text{con-private-coins}, *)$

- “*”: ‘expected budget’ — expected consumption should be at most this much

Definitions and notation

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Main result

$$\mathcal{IP}_{\text{pub-ran}}(\text{con-space, con-private-coins}, *) = \text{P}$$

- $\mathcal{IP}_{\text{pub-ran}}(\text{log-space}) = \text{P}$ [Con89; GKR15]
- $\mathcal{IP}_{\text{pub-ran}}(\text{log-space}) = \mathcal{IP}_{\text{pub-ran}}(\text{con-space, con-private-coins}, *)$
 - A) $\mathcal{IP}_{\text{pub-ran}}(\text{log-space}) \subseteq \mathcal{IP}_{\text{pub-ran}}(\text{con-space, con-private-coins}, *)$
 - B) $\mathcal{IP}_{\text{pub-ran}}(\text{con-space, con-private-coins}, *) \subseteq \mathcal{IP}_{\text{pub-ran}}(\text{log-space})$

Main result

$$\text{A) } \mathcal{IP}_{\text{pub-ran}}(\text{log-space}) \subseteq \mathcal{IP}_{\text{pub-ran}}(\text{con-space, con-private-coins}, *)$$

$$1 \quad \mathcal{IP}_{\text{pub-ran}}(\text{log-space}) \text{ to } \mathcal{IP}_{\text{pub-ran}}(\text{con-space}, k\text{-head})$$

[...]

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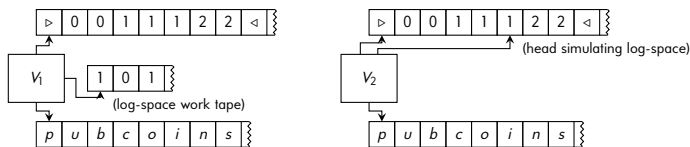
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1 $\mathcal{IP}_{\text{pub-ran}}(\text{log-space})$ to $\mathcal{IP}_{\text{pub-ran}}(\text{con-space}, k\text{-head})$

- Direct simulation via encoding $\log(n)$ -bit memory into head positions [Har72]
- $O(n)$ time overhead



2 $\mathcal{IP}_{\text{pub-ran}}(\text{con-space}, k\text{-head})$ to $\mathcal{IP}_{\text{pub-ran}}(\text{con-space, con-private-coins}, *)$

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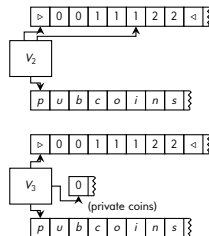
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■ Simulate

- 1 Pick head i with pr. p_i .
- 2 Simulate the 'story' of k heads, tracing and verifying i th head.
- 3 Interactivity is passed on.



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■ Repeat for m rounds:

■ Simulate

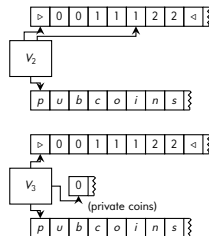
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■ Simulate with pr. $p_{\text{simulation}} = \sum_i p_i$.

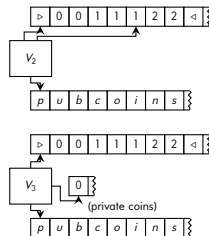
1 Pick head i with pr. p_i .

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3 Interactivity is passed on.

■ Time the 'story' with the remaining high pr., p_{timer} . [GS22]

■
■
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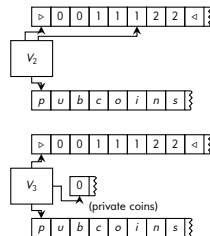
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■ Time the 'story' with the remaining high pr., p_{timer} . [GS22]

1 Random walk t times.

2 Abort and *reject* if all end on $\triangleleft$. Else, repeat.

■ Markov analysis: $O(n^{t+2})$ expected; almost always more than $O(n^{t+1})$



Main result

B) $\mathcal{IP}_{\text{pub-ran}}(\text{con-space}, \text{con-private-coins}, *) \subseteq \mathcal{IP}_{\text{pub-ran}}(\text{log-space})$

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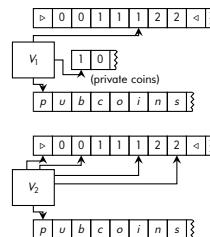
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- r random bits $\leadsto 2^r$ heads to parallelly simulate 2^r probable executions
 - 2^r interactions multiplexed
 - Verify: sims that interacted identically are responded the same
- t more heads to limit runtime to $O(n^t)$, deterministically
- Decide by a randomly chosen sim

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[...]



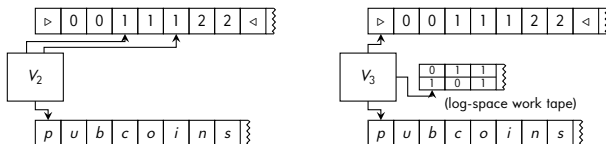
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- Direct simulation by keeping head indices in $\log(n)$ -bit memory [Har72]
- $O(n)$ time overhead, can be improved to $O(n/\log n)$



Main result (precise)

$$P = \text{IP}(\text{log-space, poly-public-coins, poly-time}) = \\ \text{IP}(\text{con-space, con-private-coins, poly}^*\text{-public-coins, poly}^*\text{-time})$$

For any $t > 1$:

- A) $\text{IP}(\text{log-space, } O(n^t)\text{-public-coins, } O(n^t)\text{-time}) \subseteq \\ \text{IP}(\text{con-space, con-private-coins, } O(n^{t+2})^*\text{-public-coins, } O(n^{t+2})^*\text{-time})$
- B) $\text{IP}(\text{con-space, con-private-coins, } O(n^t)^*\text{-public-coins, } O(n^t)^*\text{-time}) \subseteq \\ \text{IP}(\text{log-space, } O(n^{\lceil t \rceil + 1})\text{-public-coins, } O(n^{\lceil t \rceil + 1})\text{-time})$

Side result

$$\text{NC} \subseteq \text{IP}(\text{con-space}, \text{con-private-coins}, O(n^4)^* \text{-public-coins}, O(n^4)^* \text{-time})$$

- $\text{NC} \subseteq \text{IP}(\text{log-space}, O(n \log^2 n) \text{-public-coins}, O(n \log^2 n) \text{-time})$ [FL93]
 - Thus, $\text{NC} \subseteq \text{IP}(\text{log-space}, O(n^2) \text{-public-coins}, O(n^2) \text{-time})$
- Result follows from “A)”

References

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- [Har72] Juris Hartmanis. “On non-determinacy in simple computing devices”. en. In: *Acta Informatica* 1.4 (1972), pp. 336–344.